

Uncertainty, vagueness and probability of many-valued events

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- 1 Why fuzzy logic?
- 2 Uncertainty vs vagueness
- 3 Probability of many-valued events

Motivation

Classical logic is the main tool for formalizing reasoning, **but**

- its expressive power is not enough to formalize many facets of commonsense reasoning;
- there is a need to cope with different forms of imperfect information: partial, uncertain, imprecise, vague, etc.

Mathematical Fuzzy Logic

[Hájek, 1998]

- formal systems (syntax, semantics, complete axiomatizations, proof theory, etc...)
- $[0, 1]$: usual choice of truth-value set
- truth-functionality assumption
- logics of comparative truth: $\text{truth}(\phi \rightarrow \psi) = 1$ iff $\text{truth}(\phi) \leq \text{truth}(\psi)$
- generalizations of classical logic

Mathematical Fuzzy Logic

Hàjek's idea: to base the semantics on the truth function for conjunction:

A **t-norm** is a binary operation $*$ on $[0, 1]$ such that:

- (i) $*$ is commutative and associative,
- (ii) $*$ is non-decreasing in both arguments,
- (iii) $1 * x = x$ and $0 * x = 0$ for all $x \in [0, 1]$.

The choice of the t-norm determines the **whole** calculus, indeed the truth function of implication is the residuum of the t-norm:

$$x \rightarrow y = \sup\{z : x * z \leq y\}$$

(if the t-norm is continuous, such *sup* exists and it is unique)

Three main logics:

- Łukasiewicz logic Ł ['20s - '30s]
 - $x *_Ł y = \max(0, x + y - 1)$
 - $x \rightarrow_Ł y = \min(1, 1 - x + y)$
 - $\neg_Ł x = x \rightarrow_Ł 0 = 1 - x$
- Gödel logic G [1930 Heyting, 1933 Gödel, 1959 Dummett]
 - $x *_G y = \min(x, y)$
 - $x \rightarrow_G y = 1$ if $x \leq y$, or $x \rightarrow_G y = y$ otherwise
 - $\neg_G x = 1$ if $x = 0$, or $\neg_G x = 0$ otherwise
- Product logic Π [Esteva, Godo, Hájek 1996]
 - $x *_G y = x \cdot y$
 - $x \rightarrow_Π y = 1$ if $x \leq y$, or $x \rightarrow_Π y = y/x$ otherwise
 - $\neg_Π x = 1$ if $x = 0$, or $\neg_Π x = 0$ otherwise

Why Ł, G and II?

- Hájek's framework is well-established and deeply studied. Between fuzzy logics given by continuous t-norms, Ł, G and II are **fundamental**: any other such logic is a combination of them.
- They enjoy interesting and useful properties. For example, the algebra on $[0, 1]$ is **standard**: the algebra of formulas with n variables corresponds exactly to the algebra of $[0, 1]$ -valued functions with domain $[0, 1]^n$ and operations defined componentwise by standard ones.

$$\phi \longleftrightarrow f_\phi$$

with ϕ formula of n variables, $f_\phi : [0, 1]^n \rightarrow [0, 1]$.

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- UNCERTAINTY \implies *PROBABILITY*

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In particular:

- *Many-valued logics* deal with **vague** concepts and they use **intermediate truth values**,
- *Probability* deals with events that are **uncertain** now, but that will become true or false later, and it uses **degrees of belief**.

Uncertainty vs vagueness

- Think of a drink that is poisonous with truth-degree 0.1 or a drink with probability $1/10$ to be poisonous.



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- Fuzzy logics are truth functional:

$$\text{truth}(A \& B) = \text{truth}(A) \& \text{truth}(B)$$

while probability is not:

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Probability of vague events

A connection: what does it mean to speak about probability of many-valued events?

Will there be traffic?, is it going to be cold tonight?

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(CLASSICAL) PROBABILITY THEORY \implies STATE THEORY

Classical probability functions

Let X be a nonempty set of events. Let \mathcal{B} a collection of subsets of X , closed by intersection, union and complement, containing \emptyset and X (i.e. a *Boolean algebra*).

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A **finitely additive probability** is a function $P : \mathcal{B} \rightarrow [0, 1]$ such that:

(i) If $A, B \in \mathcal{B}$, where $A \cap B = \emptyset$, then

$$P(A \cup B) = P(A) + P(B),$$

(ii) $P(\emptyset) = 0$ and $P(X) = 1$.

Probability of many-valued events: states

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A **state** is a map $s : A \rightarrow [0, 1]$ such that:

(i) For every $a, b \in A$, if $a *_L b = 0$, then

$$s(a +_L b) = s(a) + s(b),$$

(ii) $s(1) = 1$.

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(ii) $s(1) = 1$.

The condition means additivity with respect to Łukasiewicz sum $+_L$.

Thus states can be thought of as generalizations of finitely additive probabilities.

States: a developing theory

- States of **Łukasiewicz** logic
D. Mundici, *Averaging the Truth-value in Łukasiewicz Logic*. *Studia Logica* 55(1), **1995**.
- States of **Gödel** logic
S. Aguzzoli, B. Gerla, V. Marra, *Defuzzifying formulas in Gödel logic through finitely additive measures*. *Proceedings FUZZ-IEEE*, **2008**.
- States of **product** logic
L. Godo, T. Flaminio, S. Ugolini *States of free product algebras and their integral representation*, to appear

States: why are they relevant?

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- States of \mathbb{L} , G and Π are connected to regular Borel probability measures. This allows to regard them as **expected values** of bounded random variables.
- States can be regarded as operators **averaging the truth value** of \mathbb{L} , G , Π logics.
- States characterize the *coherence* criterion of de Finetti's foundation of subjective probability wrt many-valued events. In this sense, states are **subjective probability measures**.

Integral representation of states

Let \mathbf{A} be the algebra of formulas of n variables of \mathbf{L} , \mathbf{G} or \mathbf{II} respectively.

A map $s : A \rightarrow [0, 1]$ is a **state** iff there is a unique (regular Borel) **probability measure** μ over $[0, 1]^n$ such that, for every $f_\phi \in A$,

$$s(f_\phi) = \int_{[0,1]^n} f_\phi \, d\mu.$$

Łukasiewicz: Kroupa (2005) - Panti (2009)

Gödel: Aguzzoli, Gerla, Marra (2008)

Product: Godo, Flaminio, U. (2017)

Expected value

Let X be a finite set and let $A = [0, 1]^X$ be the algebra of Łukasiewicz functions from X in $[0, 1]$.

Every $f \in A$ can be regarded as a **real-valued and bounded random variable** on X .

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Thus, via the integral representation, states can be seen as **expected values** of f , indeed:

$$E(f) = \int_X f_\phi \, d\mu = s_\mu(f)$$

Averaging the truth value

The integral representation allows us to associate a **real value** to each formula of the logic:

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Moreover, for all \mathfrak{L} , \mathcal{G} and Π , we can prove that each possible state belongs to the **convex closure** of the **valuations of the logic**.

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Frequentist answer: build a huge number of bridges, wait for 200 years and compute the ratio between the number of bridges which resisted and the total number of bridges.

de Finetti's foundation of subjective probability: coherent betting odds

Events of interest: $e_1 \dots e_k$.

Bookmaker publishes a book β assigning a betting odd $\beta_i \in [0, 1]$ to each e_i .

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Bookmaker pays back to the gambler σ_i euros if e_i turns out to be true in w , or nothing if it is false in w .

Total balance: $\sum_{i=1}^k \sigma_i (\beta_i - w(e_i))$.

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Example: How much would you bet on the bridge between Reggio and Messina resisting for 200 years?

de Finetti's foundation of subjective probability: coherent betting odds

A suitable formalization of classical de Finetti's betting game consists in interpreting events, books and possible worlds this way:

- 1 **events** are elements of an arbitrary boolean algebra \mathbf{B} ,
- 2 a **book** on a finite subset $\{e_1, \dots, e_k\} \subseteq B$ is a map $\beta : e_i \mapsto \beta_i \in [0, 1]$,
- 3 a **possible world** is a structure preserving map from \mathbf{B} into the two element boolean algebra $\mathbf{2}$, that is, any element of $\mathcal{H}(\mathbf{B}, \mathbf{2})$.

de Finetti's foundation of subjective probability

Classical Coherence Criterion

Let \mathbf{B} be a boolean algebra and let $\{e_1, \dots, e_k\}$ be a finite subset of B . A book $\beta : e_i \mapsto \beta_i$ is said to be **coherent** iff for each choice of $\sigma_1, \dots, \sigma_k \in \mathbb{R}$, there exists $w \in \mathcal{H}(\mathbf{B}, \mathbf{2})$ such that:

$$\sum_{i=1}^k \sigma_i (\beta_i - w(e_i)) \geq 0$$

Theorem

Let \mathbf{B} be a boolean algebra, $B' = \{e_1, \dots, e_k\}$ be a finite subset of B and let β be a book on B' . Then the following are equivalent:

- 1 β is *coherent*.
- 2 There exists a *probability* p of \mathbf{B} such that p coincides with β over B' .

de Finetti's foundation of subjective probability

De Finetti never considered the case of many-valued events, anyway it is not difficult to reframe his coherence criterion in the many-valued realm:

Many-valued Coherence Criterion

Let \mathbf{A} be an MV-algebra and $A' = e_1, \dots, e_k$ be a finite subset of A . We say that a book $\beta : e_i \mapsto \beta_i$ is **coherent** iff for each choice of $\sigma_1, \dots, \sigma_k \in \mathbb{R}$, there exists $w \in \mathcal{H}(\mathbf{A}, [0, 1]_{MV})$ such that

$$\sum_{i=1}^k \sigma_i (\beta_i - w(e_i)) \geq 0.$$

Theorem

Let \mathbf{A} be an algebra of Łukasiewicz logic, $A' = \{e_1, \dots, e_k\}$ be a finite subset of A and let β be a book on A' . Then the following are equivalent:

- ① β is *coherent*.
- ② There exists a *state* s of \mathbf{A} such that s coincides with β over A' .

Recap

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- States as operators averaging the truth value of \mathbb{L} , G , Π logics;
- States as subjective probability measures in de Finetti's theory.